

Daisy Summary

These figures and indices are assembled from the the full manual:

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1A. Daisy Expressions

Expressions, E numerals, N literals, I $\boxed{\text{"}} - characters - \boxed{\text{"}}$ (literal quotation) $\boxed{\text{`}} E \boxed{\text{'}}$ (value quotation) $\boxed{\text{[}} E \boxed{\text{]}}$ (parentheses) $E_0 \boxed{\text{:}} E_1$ (application expression) $\boxed{\text{[}} X \boxed{\text{]}} E$ (function expression) $\boxed{\text{[}} \boxed{\text{]}}$ (Nil) $\boxed{\text{[}} E_0 \cdots E_n \boxed{\text{]}}$ (list value) $\boxed{\text{[}} E_0 \cdots E_n \boxed{\text{!}} E_{n+1} \boxed{\text{]}}$ (list value) $\boxed{\text{[}} E_0 \cdots E_n \boxed{\text{*}} \boxed{\text{]}}$ (list value) $\boxed{\text{<}} \boxed{\text{>}}$ (Nil) $\boxed{\text{<}} E_0 \cdots E_n \boxed{\text{>}}$ (list expression) $\boxed{\text{<}} E_0 \cdots E_n \boxed{\text{!}} E_{n+1} \boxed{\text{>}}$ (list expression) $\boxed{\text{<}} E_0 \cdots E_n \boxed{\text{*}} \boxed{\text{>}}$ (list expression) $\boxed{\text{[}} \boxed{\text{]}}$ (Nil) $\boxed{\text{[}} E_0 \cdots E_n \boxed{\text{]}}$ (list expression) $\boxed{\text{[}} E_0 \cdots E_n \boxed{\text{!}} E_{n+1} \boxed{\text{]}}$ (list expression) $\boxed{\text{[}} E_0 \cdots E_n \boxed{\text{*}} \boxed{\text{]}}$ (list expression)**Formal Arguments, X**

literals

 $\boxed{\text{[}} \boxed{\text{]}}$ $\boxed{\text{[}} X_0 \cdots X_n \boxed{\text{]}}$ $\boxed{\text{[}} [X_0 \cdots X_n \boxed{\text{!}} X_{n+1} \boxed{\text{]}}$

1B. Composite Forms

Binding Forms

```
let:[ X E0 E1 ]
ret:[ X E0 E1 ]
fix:[ X ! E1 ]
```

Conditional Form

```
if:< E0 E'0 … En E'n En+1 >
```

Surface Assignment

literal \equiv E

1C. A Daisy Grammar (plus surface assignment)

Atomic Expressions

EXP ::= NUMERAL | LITERAL

List Expressions

*EXP ::= [ELTS] | { ELTS } | [] ELTS |
 ELTS ::= —nothing— | EXP ! EXP | EXP * | EXP ELTS*

Application Expression

EXP ::= EXP [] EXP

Function Expression

*EXP ::= [] ARG [.] EXP
 ARG ::= LITERAL | [] ARGS []
 ARGS ::= —nothing— | ARG ARGS | ARG [!] ARG*

Quotation Expressions

*EXP ::= [" —characters— "]
 EXP ::= [~] Exp*

Parenthesized Expression

EXP ::= [() EXP]

Binding Forms

let: [ARG EXP EXP]
rec: [ARG EXP EXP]
fix: [ARG ! EXP]

Conditional Form

if: < EXP Exp ⋯ EXP >

Surface Assignment

LITERAL [=] EXP

2A. Evaluation

$\mathcal{V}_\rho[E]$ is the value of expression E in environment ρ . Variables E stand for expressions, I for literal identifiers, and X for function arguments. Lower case variables stand for values.

$$\mathcal{V}_\rho[N] = N$$

$$\mathcal{V}_\rho[I] = \begin{cases} v, & \text{if } v \text{ is assigned to } I; \\ \rho(I), & \text{otherwise.} \end{cases}$$

$$\mathcal{V}_\rho[\text{add}] = \text{add}$$

$$\mathcal{V}_\rho["—characters—"] = \text{the literal } —characters—$$

$$\mathcal{V}_\rho[\wedge E] = E$$

$$\mathcal{V}_\rho[(E)] = \mathcal{V}_\rho[E]$$

$$\mathcal{V}_\rho[\lambda X . E] = {}_\rho\lambda X . E$$

$$\mathcal{V}_\rho[E_0 : E_1] = \mathcal{A}_\rho(\mathcal{V}_\rho[E_0])(\mathcal{V}_\rho[E_1])$$

$$\mathcal{V}_\rho[[\]]=[\]$$

$$\mathcal{V}_\rho[[E_0 \dots E_n]] = [E_0 \dots E_n]$$

$$\mathcal{V}_\rho[[E_0 \dots E_n ! E_{n+1}]] = [E_0 \dots E_n ! E_{n+1}]$$

$$\mathcal{V}_\rho[[E_0 \dots E_n *]] = [E_0 \dots E_n *]$$

$$\mathcal{V}_\rho[<>] = [\]$$

$$\mathcal{V}_\rho[<E_0 \dots E_n>] = [\mathcal{V}_\rho[E_0] \dots \mathcal{V}_\rho[E_n]]$$

$$\mathcal{V}_\rho[<E_0 \dots E_n ! E_{n+1}>] = [\mathcal{V}_\rho[E_0] \dots \mathcal{V}_\rho[E_n] ! \mathcal{V}_\rho[E_{n+1}]]$$

$$\mathcal{V}_\rho[<E_0 \dots E_n *>] = [\mathcal{V}_\rho[E_0] \dots \mathcal{V}_\rho[E_n]*]$$

2B. Application

$\mathcal{A}_\rho u v$ interprets the application of value u to value v in environment ρ . Variables E stand for expressions, X for formal arguments, N for numerals.

$$\mathcal{A}_\rho \text{add} [u \ v] = u + v$$

$$\mathcal{A}_\rho N [v_0 \ v_1 \ \dots] = v_N$$

$$\mathcal{A}_\rho (\rho' \setminus X . E) v = \mathcal{V}_{\rho''} [E] \text{ where } \rho'' = \rho' \left[\begin{matrix} v \\ X \end{matrix} \right]$$

$$\mathcal{A}_\rho \text{val} v = \mathcal{V}_\rho [v]$$

$$\mathcal{A}_\rho \text{let} [X \ E \ E'] = \mathcal{V}_{\rho'} [E'] \text{ where } \begin{cases} \rho' = \rho \left[\begin{matrix} v \\ X \end{matrix} \right] \\ v = \mathcal{V}_\rho [E_0] \end{cases}$$

$$\mathcal{A}_\rho \text{rec} [X \ E \ E'] = \mathcal{V}_{\rho'} [E'] \text{ where } \begin{cases} \rho' = \rho \left[\begin{matrix} v \\ X \end{matrix} \right] \\ v = \mathcal{V}_{\rho'} [E_0] \end{cases}$$

$$\mathcal{A}_\rho \text{fix} [X ! E] = v \text{ where } \begin{cases} \rho' = \rho \left[\begin{matrix} v \\ X \end{matrix} \right] \\ v = \mathcal{V}_{\rho'} [E_0] \end{cases}$$

$$\mathcal{A}_\rho \text{if} [u_0 \ v_0 \ \dots \ u_n \ v_n \ w] = \begin{cases} v_i, & \text{if } u_i \neq [] \text{ and } u_0 \dots u_{i-1} = [] \\ w, & \text{if } u_0 \dots u_n = [] \end{cases}$$

$$\mathcal{A}_\rho [] (v) = []$$

$$\mathcal{A}_\rho [v ! v'] [[u_0 ! u'_0] \ \dots] = [w ! w'] \text{ where } \begin{cases} w = \mathcal{A}_\rho v [u_0 \ \dots] \\ w' = \mathcal{A}_\rho v' [u'_0 \ \dots] \end{cases}$$

$$\mathcal{A}_\rho [v *] [[\] \ \dots] = [\]$$

$$\mathcal{A}_\rho [v *] [[u_0 ! u'_0] \ \dots] = [w ! w'] \text{ where } \begin{cases} w = \mathcal{A}_\rho v [u_0 \ \dots] \\ w' = \mathcal{A}_\rho [v *] [u'_0 \ \dots] \end{cases}$$

2C. Other Values

$$\mathcal{A}_\rho \perp v = \perp$$

$$\mathcal{A}_\rho v \perp = \perp$$

$$\mathcal{A}_\rho \text{erron } v = \text{erron}$$

$$\mathcal{A}_\rho [I] v = \text{error}$$

$$\mathcal{A}_\rho [E_0 : E_1] v = \text{error}$$

$$\mathcal{A}_\rho [\lambda X . E] v = \mathcal{V}_\rho [E] \text{ where } \rho = \phi \begin{bmatrix} v \\ X \end{bmatrix}$$

$$\mathcal{V}_\rho \perp = \perp$$

$$\mathcal{V}_\rho \text{erron} = \text{erron}$$

$$\mathcal{V}_\rho (\rho \lambda X . E) = \text{illformed closure}$$

2D. Indeterminate Lists

The indices $0', 1', \dots, n'$, $n \geq 0$, are a rearrangement of $0, 1, \dots, n$. Variables $A, M, \Omega, E_0, E_1, \dots$ stand for expressions whose values are $a, m, \perp, e_0, e_1, \dots$, respectively. Assume a is a non-list, m is an error, and that each of the e_i s is valid.

$$\mathcal{V}_\rho[\{\}] = []$$

$$\mathcal{V}_\rho[\{E_0 \dots E_n\}] \stackrel{\Delta}{=} [e_{0'} \dots e_{n'}]$$

$$\mathcal{V}_\rho[\{E_0 ! \{E_1 \dots E_n\}\}] \stackrel{\Delta}{=} [e_{0'} \dots e_{n'}]$$

$$\mathcal{V}_\rho[\{\Omega \ E_0 \dots E_n\}] \stackrel{\Delta}{=} [e_{0'} \dots e_{n'} ! \perp]$$

$$\mathcal{V}_\rho[\{M \ E_0 \dots E_n\}] \stackrel{\Delta}{=} [e_{0'} \dots e_{n'} \ m]$$

$$\mathcal{V}_\rho[\{E_0 ! \{E_1 ! E_2\}\}] \stackrel{\Delta}{=} \mathcal{V}_\rho[\{E_0 \ E_1 ! E_2\}]$$

$$\mathcal{V}_\rho[\{E_0 ! \langle E_1 ! E_2 \rangle\}] \stackrel{\Delta}{=} \begin{cases} [e_0 ! \mathcal{V}_\rho[\langle E_1 ! E_2 \rangle]] \\ [e_1 ! \mathcal{V}_\rho[\{E_0 ! E_2\}]] \end{cases}$$

$$\mathcal{V}_\rho[\{E_0 \dots E_n \ E_{n+1} *\}] \stackrel{\Delta}{=} [e_{0'} \dots e_{n'} \ e_{n+1} *]$$

$$\mathcal{V}_\rho[\{E_0 \dots E_n ! A\}] \stackrel{\Delta}{=} [e_{0'} \dots e_{n'} ! a]$$

3. Environments

Environments ρ associate values with literal identifiers. $\rho(I)$ is I 's binding in ρ . Environment extension is specified as

$$\begin{aligned}\rho\left[\begin{matrix} v \\ I \end{matrix}\right](J) &= \begin{cases} v, & \text{if } I = J \\ \rho(J), & \text{if } I \neq J \end{cases} \\ \rho\left[\begin{matrix} [u ! v] \\ [X ! Y] \end{matrix}\right](J) &= \rho\left[\begin{matrix} v \\ Y \end{matrix}\right]\left[\begin{matrix} v \\ X \end{matrix}\right](J) \\ \rho\left[\begin{matrix} v \\ [] \end{matrix}\right](J) &= \rho(J)\end{aligned}$$

Environments accept undefined bindings,

$$\rho\left[\begin{matrix} \perp \\ I \end{matrix}\right](J) = \begin{cases} \perp, & \text{if } I = J \\ \rho(J), & \text{if } I \neq J \end{cases}$$

The void environment, \star binds every literal to an unbound-identifier error,

$$\star(I) = |\text{ubi}:I|$$

4. Closure Objects

The closure of a function expression $\lambda X . E$ is denoted $\rho\lambda X . E$, and appears in output as $\lambda?\lambda X . E$.

Daisy Operations (by group)

Types in descriptors		
\mathcal{B} – T or []		
\mathcal{N} – a numeral	\mathcal{L} – a list	\mathcal{C} – a character
\mathcal{I} – a literal	\mathcal{A} – an atom	\mathcal{E} – an expression
\mathcal{V} – a value		

Arithmetic Tests

<code>zero?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	<i>Test for Zero</i>
<code>one?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	<i>Test for One</i>
<code>neg?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	<i>Test for negative</i>
<code>pos?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	<i>Test for positive</i>
<code>lt?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$</code>	<i>Less-than</i>
<code>le?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$</code>	<i>At-most</i>
<code>eq?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$</code>	<i>Numeric equality</i>
<code>ne?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$</code>	<i>Numeric inequality</i>
<code>ge?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$</code>	<i>At-least</i>
<code>gt?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$</code>	<i>Greater-than</i>

Unary Arithmetic

<code>neg: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$</code>	<i>Negate</i>
<code>sgn: $\mathcal{N} \rightarrow -1 \text{ or } 1$</code>	<i>Sign projection</i>
<code>inc: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$</code>	<i>Increment</i>
<code>dcr: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$</code>	<i>Decrement</i>
<code>inv: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$</code>	<i>Invert</i>

Binary Arithmetic

add: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Add
sub: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Subtract
mpy: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Multiply
div: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Divide
rem: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Remainder

Binary Logic

or: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Logical-or
and: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Logical-and
xor: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Exclusive-or

Reference Tests

nil?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for Nil
isNML?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a numeral
isLltr?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a literal
isAtm?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for an Atom
isLST?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a list
same?: $[\mathcal{U} \ \mathcal{V}_1 \dots \mathcal{V}_n] \rightarrow \mathcal{B}$	Reference equality

List Processing

head: $[\mathcal{V}_1 ! \mathcal{V}_2] \rightarrow \mathcal{V}_1$	Head of a list
tail: $[\mathcal{V}_1 ! \mathcal{V}_2] \rightarrow \mathcal{V}_1$	Tail of a list
cons: $[\mathcal{V}_1 \ \mathcal{V}_2] \rightarrow [\mathcal{V}_1 ! \mathcal{V}_2]$	List constructor
frons: $[\mathcal{V} \ \mathcal{V}'] \rightarrow \mathcal{L}$	Multiset constructor
any?: $[\mathcal{V}_0 \ \mathcal{V}_1 \dots] \rightarrow [\mathcal{V}_i \ \mathcal{V}_{i+1} \dots]$	Locate a non-Nil element
all?: $[\mathcal{V}_0 \ \mathcal{V}_1 \dots] \rightarrow \mathcal{B}$	Test for no null elements
in?: $[\mathcal{U} \ [\mathcal{V}_0 \ \mathcal{V}_1 \dots]] \rightarrow \mathcal{B}$	Membership Test
if: $[\mathcal{T}_0 \ \mathcal{V}_0 \ \mathcal{T}_1 \ \mathcal{V}_1 \dots] \rightarrow \mathcal{V}_i$	Conditional operation

Object Manipulation

TagOf: $\mathcal{V} \rightarrow \mathcal{N}$	Numeric value of a tag
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Tag Tests

isDCT?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a directive
isNML?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a numeral
isLST?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a list-object
isAPL?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a application-object
isFTN?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a function-object
isIDE?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a identifier-object
isERR?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for an identifier-object

Tag Coercion

asDCT: $\mathcal{V} \rightarrow \mathcal{D}$	Cite as a directive
asNML: $\mathcal{V} \rightarrow \mathcal{N}$	Cite as a numeral
asIDE: $\mathcal{V} \rightarrow \mathcal{I}$	Cite as an identifier
asLST: $\mathcal{V} \rightarrow \mathcal{L}$	Cite as a list
asAPL: $\mathcal{V} \rightarrow \mathcal{A}$	Cite as an application
asFTN: $\mathcal{V} \rightarrow \mathcal{F}$	Cite as a function
asERR: $\mathcal{V} \rightarrow \mathcal{M}$	Error

Access to Composite Objects

_hd: $\mathcal{V} \rightarrow \mathcal{V}'$	Head of a composite object
_tl: $\mathcal{V} \rightarrow \mathcal{V}'$	Tail of a composite object

Character Manipulation

Chr?: $\mathcal{V} \rightarrow \mathcal{B}$	Character test
ChrAsNml: $\mathcal{C} \rightarrow \mathcal{N}$	Character's numeric code
NmlAsChr: $\mathcal{N} \rightarrow \mathcal{C}$	Convert a numeral to a character

Character Classification

ScnSPC?: $\mathcal{C} \rightarrow \mathcal{B}$	Space-character test
ScnDGT?: $\mathcal{C} \rightarrow \mathcal{B}$	Digit-character test
ScnLFA?: $\mathcal{C} \rightarrow \mathcal{B}$	Alpha-character test
ScnNON?: $\mathcal{C} \rightarrow \mathcal{B}$	Neutral-character test
ScnSYM?: $\mathcal{C} \rightarrow \mathcal{B}$	Symbol-character test
ScnCTL?: $\mathcal{C} \rightarrow \mathcal{B}$	Control-character test

Sequencing

<code>crc_hd: $\mathcal{V} \rightarrow \mathcal{V}$</code>	<i>Coerce-head</i>
<code>crc_tl: $\mathcal{V} \rightarrow \mathcal{V}$</code>	<i>Coerce-tail</i>
<code>seq: $[\mathcal{V}_0 \ \mathcal{V}_1 \ \dots \ \mathcal{V}_n] \rightarrow \mathcal{V}_n$</code>	<i>Sequencer</i>

Interface Operations

<code>console: $\mathcal{I} \rightarrow [\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]$</code>	<i>Interactive input</i>
<code>screen: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow []$</code>	<i>Interactive output</i>
<code>dski: $\mathcal{I} \rightarrow [\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]$</code>	<i>File input</i>
<code>dsko: $[\mathcal{I}[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]] \rightarrow []$</code>	<i>File output</i>

Text Generation

<code>issue: $\mathcal{V} \rightarrow [\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]$</code>	<i>Generate text</i>
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Scanning

<code>scan: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow [\mathcal{A} \ \mathcal{C}_j \ \mathcal{C}_{j+1} \ \dots]$</code>	<i>Scan text</i>
<code>scans: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow [\mathcal{A}_0 \ \mathcal{A}_1 \ \dots]$</code>	<i>scan iterated</i>

Parsing

<code>parse: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow [\mathcal{E}_0 \ \mathcal{E}_1 \ \dots]$</code>	<i>xpareses \circ scans</i>
<code>xparse: $[\mathcal{T}_0 \ \mathcal{T}_1 \ \dots] \rightarrow [\mathcal{E} \ \mathcal{T}_j \ \mathcal{T}_{j+1} \ \dots]$</code>	<i>Parse text</i>
<code>xpareses: $[\mathcal{T}_0 \ \mathcal{T}_1 \ \dots] \rightarrow [\mathcal{E}_0 \ \mathcal{E}_1 \ \dots]$</code>	<i>parse iterated</i>

Special Operations

<code>val: $\mathcal{E} \rightarrow \mathcal{V}$</code>	<i>Evaluate</i>
<code>evlst: $[\mathcal{E}_0 \ \mathcal{E}_1 \ \dots] \rightarrow [\mathcal{V}_0 \ \mathcal{V}_1 \ \dots]$</code>	<i>val iterated.</i>
<code>let: $[\mathcal{X} \ \mathcal{E}_1 \ \mathcal{E}_2] \rightarrow \mathcal{V}$</code>	<i>Lexical binder</i>
<code>rec: $[\mathcal{X} \ \mathcal{E}_1 \ \mathcal{E}_2] \rightarrow \mathcal{V}$</code>	<i>Recursive binder</i>
<code>fix: $[\mathcal{X} \ \mathcal{E}] \rightarrow \mathcal{V}$</code>	<i>Recursive binder</i>

Daisy Operations (by group)

Types in descriptors
$\mathcal{B} - \mathbf{T}$ or $[]$
\mathcal{N} – a numeral \mathcal{L} – a list \mathcal{C} – a character
\mathcal{I} – a literal \mathcal{A} – an atom \mathcal{E} – an expression
\mathcal{V} – a value

add: $[\mathcal{N}_1 \ \mathcal{N}_2] \longrightarrow \mathcal{N}_3$	<i>Add</i>
all?: $[\mathcal{V}_0 \ \mathcal{V}_1 \ \dots] \longrightarrow \mathcal{B}$	<i>Test for no null elements</i>
and: $[\mathcal{N}_1 \ \mathcal{N}_2] \longrightarrow \mathcal{N}_3$	<i>Logical-and</i>
any?: $[\mathcal{V}_0 \ \mathcal{V}_1 \ \dots] \longrightarrow [\mathcal{V}_i \ \mathcal{V}_{i+1} \ \dots]$	<i>Locate a non-Nil element</i>
asAPL: $\mathcal{V} \longrightarrow \mathcal{A}$	<i>Cite as an application</i>
asDCT: $\mathcal{V} \longrightarrow \mathcal{D}$	<i>Cite as a directive</i>
asERR: $\mathcal{V} \longrightarrow \mathcal{M}$	<i>Error</i>
asFTN: $\mathcal{V} \longrightarrow \mathcal{F}$	<i>Cite as a function</i>
asIDE: $\mathcal{V} \longrightarrow \mathcal{I}$	<i>Cite as an identifier</i>
asLST: $\mathcal{V} \longrightarrow \mathcal{L}$	<i>Cite as a list</i>
asNML: $\mathcal{V} \longrightarrow \mathcal{N}$	<i>Cite as a numeral</i>
Chr?: $\mathcal{V} \longrightarrow \mathcal{B}$	<i>Character test</i>
ChrAsNml: $\mathcal{C} \longrightarrow \mathcal{N}$	<i>Character's numeric code</i>
console: $\mathcal{I} \longrightarrow [\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]$	<i>Interactive input</i>
cons: $[\mathcal{V}_1 \ \mathcal{V}_2] \longrightarrow [\mathcal{V}_1 \ ! \mathcal{V}_2]$	<i>List constructor</i>
crc_hd: $\mathcal{V} \longrightarrow \mathcal{V}$	<i>Coerce-head</i>
crc_tl: $\mathcal{V} \longrightarrow \mathcal{V}$	<i>Coerce-tail</i>
dcr: $\mathcal{N}_1 \longrightarrow \mathcal{N}_2$	<i>Decrement</i>
div: $[\mathcal{N}_1 \ \mathcal{N}_2] \longrightarrow \mathcal{N}_3$	<i>Divide</i>
dski: $\mathcal{I} \longrightarrow [\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]$	<i>File input</i>
dsko: $[\mathcal{I}[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots]] \longrightarrow []$	<i>File output</i>

eq?: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{B}$	Numeric equality
evlst: $[\mathcal{E}_0 \mathcal{E}_1 \dots] \rightarrow [\mathcal{V}_0 \mathcal{V}_1 \dots]$	val iterated.
fix: $[\mathcal{X} \mathcal{E}] \rightarrow \mathcal{V}$	Recursive binder
frons: $[\mathcal{V} \mathcal{V}'] \rightarrow \mathcal{L}$	Multiset constructor
ge?: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{B}$	At-least
gt?: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{B}$	Greater-than
head: $[\mathcal{V}_1 ! \mathcal{V}_2] \rightarrow \mathcal{V}_1$	Head of a list
_hd: $\mathcal{V} \rightarrow \mathcal{V}'$	Head of a composite object
if: $[\mathcal{T}_0 \mathcal{V}_0 \mathcal{T}_1 \mathcal{V}_1 \dots] \rightarrow \mathcal{V}_i$	Conditional operation
in?: $[\mathcal{U} [\mathcal{V}_0 \mathcal{V}_1 \dots]] \rightarrow \mathcal{B}$	Membership Test
inc: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$	Increment
inv: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$	Invert
isAPL?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a application-object
isAtm?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for an Atom
isDCT?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a directive
isERR?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for an identifier-object
isFTN?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a function-object
isIDE?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a identifier-object
isLST?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a list-object
isLST?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a list
isLtrr?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a literal
isNML?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a numeral
isNML?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for a numeral
issue: $\mathcal{V} \rightarrow [\mathcal{C}_0 \mathcal{C}_1 \dots]$	Generate text
le?: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{B}$	At-most
let: $[\mathcal{X} \mathcal{E}_1 \mathcal{E}_2] \rightarrow \mathcal{V}$	Lexical binder
lt?: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{B}$	Less-than
mpy: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{N}_3$	Multiply
ne?: $[\mathcal{N}_1 \mathcal{N}_2] \rightarrow \mathcal{B}$	Numeric inequality
neg?: $\mathcal{N} \rightarrow \mathcal{B}$	Test for negative
neg: $\mathcal{N}_1 \rightarrow \mathcal{N}_2$	Negate
nil?: $\mathcal{V} \rightarrow \mathcal{B}$	Test for Nil
NmlAsChr: $\mathcal{N} \rightarrow \mathcal{C}$	Convert a numeral to a character

<code>one?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	Test for One
<code>or: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$</code>	Logical-or
<code>parse: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow [\mathcal{E}_0 \ \mathcal{E}_1 \ \dots]$</code>	<code>xparses</code> \circ <code>scans</code>
<code>pos?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	Test for positive
<code>rec: $[\mathcal{X} \ \mathcal{E}_1 \ \mathcal{E}_2] \rightarrow \mathcal{V}$</code>	Recursive binder
<code>rem: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$</code>	Remainder
<code>same?: $[\mathcal{U} \ \mathcal{V}_1 \ \dots \ \mathcal{V}_n] \rightarrow \mathcal{B}$</code>	Reference equality
<code>scans: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow [\mathcal{A}_0 \ \mathcal{A}_1 \ \dots]$</code>	scan iterated
<code>scan: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow [\mathcal{A} \ \mathcal{C}_j \ \mathcal{C}_{j+1} \ \dots]$</code>	Scan text
<code>ScnCTL?: $\mathcal{C} \rightarrow \mathcal{B}$</code>	Control-character test
<code>ScnDGT?: $\mathcal{C} \rightarrow \mathcal{B}$</code>	Digit-character test
<code>ScnLFA?: $\mathcal{C} \rightarrow \mathcal{B}$</code>	Alpha-character test
<code>ScnNON?: $\mathcal{C} \rightarrow \mathcal{B}$</code>	Neutral-character test
<code>ScnSPC?: $\mathcal{C} \rightarrow \mathcal{B}$</code>	Space-character test
<code>ScnSYM?: $\mathcal{C} \rightarrow \mathcal{B}$</code>	Symbol-character test
<code>screen: $[\mathcal{C}_0 \ \mathcal{C}_1 \ \dots] \rightarrow []$</code>	Interactive output
<code>seq: $[\mathcal{V}_0 \ \mathcal{V}_1 \ \dots \ \mathcal{V}_n] \rightarrow \mathcal{V}_n$</code>	Sequencer
<code>sgn: $\mathcal{N} \rightarrow -1 \text{ or } 1$</code>	Sign projection
<code>sub: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$</code>	Subtract
<code>TagOf: $\mathcal{V} \rightarrow \mathcal{N}$</code>	Numeric value of a tag
<code>tail: $[\mathcal{V}_1 ! \mathcal{V}_2] \rightarrow \mathcal{V}_1$</code>	Tail of a list
<code>_tl: $\mathcal{V} \rightarrow \mathcal{V}'$</code>	Tail of a composite object
<code>val: $\mathcal{E} \rightarrow \mathcal{V}$</code>	Evaluate
<code>xor: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$</code>	Exclusive-or
<code>xparses: $[\mathcal{T}_0 \ \mathcal{T}_1 \ \dots] \rightarrow [\mathcal{E}_0 \ \mathcal{E}_1 \ \dots]$</code>	parse iterated
<code>xparse: $[\mathcal{T}_0 \ \mathcal{T}_1 \ \dots] \rightarrow [\mathcal{E} \ \mathcal{T}_j \ \mathcal{T}_{j+1} \ \dots]$</code>	Parse text
<code>zero?: $\mathcal{N} \rightarrow \mathcal{B}$</code>	Test for Zero

arithmetic operations

Arithmetic Operations

add: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	add	$\mathcal{N}_3 = \mathcal{N}_1 + \mathcal{N}_2$
sub: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	subtract	$\mathcal{N}_3 = \mathcal{N}_1 - \mathcal{N}_2$
div: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	divide	$\mathcal{N}_3 = \mathcal{N}_1 \div \mathcal{N}_2$
mpy: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	multiply	$\mathcal{N}_3 = \mathcal{N}_1 \cdot \mathcal{N}_2$
rem: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	remainder	$\mathcal{N}_3 = \mathcal{N}_1 - (\mathcal{N}_1 \div \mathcal{N}_2) \cdot \mathcal{N}_2$
inc: $\mathcal{N} \rightarrow \mathcal{N}'$	increment	$\mathcal{N}' = \mathcal{N} + 1$
dcr: $\mathcal{N} \rightarrow \mathcal{N}'$	decrement	$\mathcal{N}' = \mathcal{N} - 1$
neg: $\mathcal{N} \rightarrow \mathcal{N}'$	negate	$\mathcal{N}' = -\mathcal{N}$
inv: $\mathcal{N} \rightarrow \mathcal{N}'$	invert	$\mathcal{N}' = \overline{\mathcal{N}}$
sgn: $\mathcal{N} \rightarrow \mathcal{N}'$	sign	$\mathcal{N}' = \begin{cases} 1, & \text{if } \mathcal{N} \geq 0 \\ -1, & \text{if } \mathcal{N} < 0 \end{cases}$
and: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	binary-and	$\mathcal{N}_3 = \mathcal{N}_1 \odot \mathcal{N}_2$
or: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	binary-or	$\mathcal{N}_3 = \mathcal{N}_1 \oplus \mathcal{N}_2$
xor: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{N}_3$	exclusive-or	$\mathcal{N}_3 = \mathcal{N}_1 \otimes \mathcal{N}_2$
zero?: $\mathcal{N} \rightarrow \mathcal{B}$	is-zero?	$\mathcal{N} = 0 ?$
one?: $\mathcal{N} \rightarrow \mathcal{B}$	is-one?	$\mathcal{N} = 1 ?$
neg?: $\mathcal{N} \rightarrow \mathcal{B}$	negative?	$\mathcal{N} < 0 ?$
pos?: $\mathcal{N} \rightarrow \mathcal{B}$	positive?	$\mathcal{N} \geq 0 ?$
lt?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	less-than	$\mathcal{N}_1 < \mathcal{N}_2 ?$
le?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	at-most	$\mathcal{N}_1 \leq \mathcal{N}_2 ?$
eq?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	equal	$\mathcal{N}_1 = \mathcal{N}_2 ?$
ne?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	unequal	$\mathcal{N}_1 \neq \mathcal{N}_2 ?$
ge?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	at-most	$\mathcal{N}_1 \geq \mathcal{N}_2 ?$
gt?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	greater-than	$\mathcal{N}_1 > \mathcal{N}_2 ?$
lt?: $[\mathcal{N}_1 \ \mathcal{N}_2] \rightarrow \mathcal{B}$	less-than	$\mathcal{N}_1 < \mathcal{N}_2 ?$

Errons

<code> —any/— </code>	<i>Erroneous list (all?).</i>
<code> —any/— </code>	<i>Erroneous list (any?).</i>
<code> —arg/— </code>	<i>Invalid formal argument.</i>
<code> —chr/— </code>	<i>Non-character operand.</i>
<code> —cmp/— </code>	<i>Erroneous comparison (same? or in?).</i>
<code> —crc/— </code>	<i>Invalid coercion.</i>
<code> —dfn/— </code>	<i>Invalid assignment.</i>
<code> —dvc/— </code>	<i>Device error.</i>
<code> —evl/— </code>	<i>Non-list argument (evlst).</i>
<code> —f-c/— </code>	<i>Construction error</i>
<code> —ftn/— </code>	<i>Error applied.</i>
<code> —hd?/— </code>	<i>Invalid head-access</i>
<code> —ifA/— </code>	<i>Invalid alternative (if).</i>
<code> —ifP/— </code>	<i>Erroneous predication (if).</i>
<code> —ld?/— </code>	<i>Invalid access</i>
<code> —lst/— </code>	<i>Error in list-expression.</i>
<code> —opn/— </code>	<i>Invalid operation code.</i>
<code> —nla/— </code>	<i>Non-list operand (head or tail).</i>
<code> —nn0/— </code>	<i>Non-numeric operand.</i>
<code> —nn1/— </code>	<i>Non-numeric operand.</i>
<code> —prb/— </code>	<i>Invalid numeric probe.</i>
<code> —sam/— </code>	<i>Invalid argument (same? or in?).</i>
<code> —scn/— </code>	<i>Invalid argument (scan, scans).</i>
<code> —seq/— </code>	<i>Invalid argument (seq).</i>
<code> —scn/— </code>	<i>Invalid argument (parse, xparses).</i>
<code> —sc0/— </code>	<i>Invalid text (scan, scans).</i>
<code> —sc1/— </code>	<i>Erroneous text (scan, scans).</i>
<code> —tag/— </code>	<i>Invalid citation.</i>
<code> —tl?/— </code>	<i>Invalid tail-access</i>
<code> —trj/— </code>	<i>Invalid trajectory.</i>
<code> —ubi:/— </code>	<i>Unbound identifier.</i>
<code> —val/— </code>	<i>Value-of-error.</i>
<code> —xps/— </code>	<i>Transposition error.</i>

Syntax Errors

Syntax Error Indicators

Indicator	Object
d	a directive
n	a numeral
m	an error
i	an identifier object
l	a list object
a	an application object
f	a function object

Parsing Messages

star is one of the indicators above.

Message	Explanation
.—	Message fill
..—	Message fill
...—	Message fill
~—	Error in value quotation
(—	Error in parenthesization
★:—	Error in application's argument
\—	Error in formal argument
\★.—	Error in function body
[—	Error before matching ‘...]’
<—	Error before matching ‘... >’
{—	Error before matching ‘... }’
!—	Error after a ‘!’
@‘=’	Error (nonliteral) before ‘=’
★=—	Error after an ‘=’

—NOTES—