

## C241 Homework Assignment 4

1. Which of the following formulas are tautologies and which are contradictions?  
Which of the formulas are logically equivalent to each other?

- (a)  $p \wedge (q \vee r)$
- (b)  $(\neg p \wedge r) \Rightarrow (q \vee r)$
- (c)  $(p \wedge q) \vee (p \vee r)$
- (d)  $\neg(r \Rightarrow q \wedge r)$
- (e)  $\neg(p \Rightarrow (q \Rightarrow p))$
- (f)  $((p \Rightarrow q) \vee (r \wedge s \vee t)) \vee (p \wedge \neg q)$

### SOLUTION

*As demonstrated by their truth tables below, (a), (c) and (d) are contingencies; (b) and (f) are tautologies; (e) is a contradiction. None of the formulas are equivalent.*

(a)

$p$	$q$	$r$	$p \wedge (q \vee r)$
$F$	$F$	$F$	$F$
$F$	$F$	$T$	$F$
$F$	$T$	$F$	$F$
$F$	$T$	$T$	$F$
$T$	$F$	$F$	$T$
$T$	$F$	$T$	$T$
$T$	$T$	$F$	$T$
$T$	$T$	$T$	$T$

(b)

$p$	$q$	$r$	$(\neg p \wedge r) \Rightarrow (q \vee r)$
$F$	$F$	$F$	$T$
$F$	$F$	$T$	$T$
$F$	$T$	$F$	$T$
$F$	$T$	$T$	$T$
$T$	$F$	$F$	$F$
$T$	$F$	$T$	$F$
$T$	$T$	$F$	$F$
$T$	$T$	$T$	$T$

(c)

$p$	$q$	$r$	$((p \wedge q) \vee (p \vee r))$
$T$	$F$	$F$	$F$
$T$	$F$	$T$	$T$
$T$	$T$	$F$	$T$
$T$	$T$	$T$	$T$
$F$	$F$	$F$	$F$
$F$	$F$	$T$	$T$
$F$	$T$	$F$	$T$
$F$	$T$	$T$	$T$



2. Use a truth table to show that

(a)  $\neg(p \vee q)$  is not logically equivalent to  $(\neg p \vee \neg q)$ , and

(b)  $\neg(p \wedge q)$  is not logically equivalent to  $(\neg p \wedge \neg q)$ .

Remember these facts.

SOLUTION

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(a)

$p$	$q$	$\neg(p \vee q)$	$\neg p \vee \neg q$
$F$	$F$	$T$	$T$
$F$	$T$	$F$	$T$
$T$	$F$	$F$	$T$
$T$	$T$	$F$	$F$

(b)

$p$	$q$	$\neg(p \wedge q)$	$\neg p \wedge \neg q$
$F$	$F$	$T$	$T$
$F$	$T$	$T$	$F$
$T$	$F$	$T$	$F$
$T$	$T$	$F$	$F$

3. Consider the logical operation defined below:

$P$	$Q$	$P \downarrow Q$
$F$	$F$	$T$
$F$	$T$	$T$
$T$	$F$	$T$
$T$	$T$	$F$

Show that ' $\downarrow$ ' can be used to implement (in the sense of Prop. 3.2) all of the operations of Definition 3.1.

SOLUTION

*By Proposition 3.2, all the operations can be implemented with ' $\vee$ ' and ' $\neg$ ', so it suffices to show how these can be implemented with ' $\downarrow$ '.*

$$\begin{array}{l}
 \neg P \rightarrow P \downarrow P \\
 P \vee Q \rightarrow (P \downarrow P) \downarrow (Q \downarrow Q) \\
 \hline
 P \wedge Q \rightarrow (P \downarrow Q) \downarrow (P \downarrow Q) \\
 P \Rightarrow Q \rightarrow P \downarrow (Q \downarrow Q) \\
 P \Leftrightarrow Q \rightarrow ((P \downarrow P) \downarrow (Q \downarrow Q)) \downarrow (P \downarrow Q) \\
 P \nabla Q \rightarrow (P \downarrow (Q \downarrow Q)) \downarrow (P \downarrow Q)
 \end{array}$$

4. Let  $P$  stand for the proposition “Sue says it.” Let  $Q$  stand for the proposition “Sam saw it.” Let  $R$  stand for the proposition “Sid did it.” Express the following sentences as formulas involving the logical connectives. If there is more than one way to translate a sentence, use truth tables to explain any differences in the meaning among these translations.

- (a) Sid did it, Sam saw it, and Sue says it.
- (b) If Sid did it, Sam saw it.
- (c) Sid did it only if Sam saw it.
- (d) Sue says it only if Sid did it, and Sam saw it.
- (e) If Sue says it implies Sam saw it, Sid did it.

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SOLUTION

- (a)  $R \wedge Q \wedge P$
- (b)  $R \Rightarrow Q$
- (c)  $R \Rightarrow Q$
- (d) *The comma indicates that the interpretation should be  $(P \Rightarrow R) \wedge Q$ . Ignoring the comma, another possibility is  $P \Rightarrow (R \wedge Q)$ . The truth table shows that these are distinct interpretations.*

$P$	$Q$	$R$	$(P \Rightarrow R) \wedge Q$	$P \Rightarrow (R \wedge Q)$
$F$	$F$	$F$	$F$	$T$
$F$	$F$	$T$	$F$	$T$
$F$	$T$	$F$	$T$	$T$
$F$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$F$	$F$
$T$	$F$	$T$	$F$	$F$
$T$	$T$	$F$	$F$	$T$
$T$	$T$	$T$	$T$	$T$

- (e)  $(P \Rightarrow Q) \Rightarrow R$ , which is not the same as  $P \Rightarrow (Q \Rightarrow R)$

$P$	$Q$	$R$	$(P \Rightarrow Q) \Rightarrow R$	$P \Rightarrow (Q \Rightarrow R)$
$F$	$F$	$F$	$F$	$T$
$F$	$F$	$T$	$T$	$T$
$F$	$T$	$F$	$F$	$T$
$F$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$T$	$T$
$T$	$F$	$T$	$T$	$T$
$T$	$T$	$F$	$F$	$F$
$T$	$T$	$T$	$T$	$T$

5. Determine whether the following proposition is a tautology.

$$(a \vee b \Leftrightarrow c) \wedge (d \vee e) \Leftrightarrow ((a \vee b \Leftrightarrow c) \wedge d) \vee ((a \vee b \Leftrightarrow c) \wedge e)$$

SOLUTION

The problem is simpler to solve if you recognize that a single term occurs three times in the proposition:

$$\boxed{(a \vee b \Leftrightarrow c)} \wedge (d \vee e) \Leftrightarrow (\boxed{(a \vee b \Leftrightarrow c)} \wedge d) \vee (\boxed{(a \vee b \Leftrightarrow c)} \wedge e)$$

Call this repeated term  $X$ . It should suffice to determine whether

$$X \wedge (d \vee e) \Leftrightarrow (X \wedge d) \vee (X \wedge e)$$

is a tautology because the original proposition can be obtained by substituting term  $\forall$  for the  $X$ . The truth table for this term has eight cases, where the original had thirty-two.

$X$	$d$	$e$	$X$	$\wedge$	$(d \vee e)$	$\Leftrightarrow$	$(X \wedge d)$	$\vee$	$(X \wedge e)$
$F$	$F$	$F$	$F$	$F$	$F$	$T$	$F$	$F$	$F$
$F$	$F$	$T$	$F$	$F$	$T$	$T$	$F$	$F$	$F$
$F$	$T$	$F$	$F$	$F$	$T$	$T$	$F$	$F$	$F$
$F$	$T$	$T$	$F$	$F$	$T$	$T$	$F$	$F$	$F$
$T$	$F$	$F$	$T$	$F$	$F$	$T$	$F$	$F$	$F$
$T$	$F$	$T$	$T$	$T$	$T$	$T$	$F$	$T$	$T$
$T$	$T$	$F$	$T$	$T$	$T$	$T$	$T$	$T$	$F$
$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$

Thus, the proposition is a tautology. By the way, we have just shown that ' $\wedge$ ' distributes over ' $\vee$ ', in the same way that multiplication distributes over addition.

In general we should check whether  $a \vee b \Leftrightarrow c$  is a tautology or a contradiction. In either case the problem is easier because the truth table is smaller. But it doesn't matter what  $X$  evaluates to;  $X \wedge (d \vee e) \Leftrightarrow (X \wedge d) \vee (X \wedge e)$  still evaluates to  $T$ .

6. Show whether the following pairs of formulas are equivalent.

(a)  $(p \Rightarrow q) \Rightarrow r$  and  $p \Rightarrow (q \Rightarrow r)$

(b)  $p \Rightarrow (q \Rightarrow r)$  and  $(p \wedge q) \Rightarrow r$

(c)  $(p \wedge q) \Rightarrow r$  and  $(p \Rightarrow r) \wedge (q \Rightarrow r)$

SOLUTION

(a)  $(p \Rightarrow q) \Rightarrow r$  and  $p \Rightarrow (q \Rightarrow r)$

*Since their truth tables differ, these formulas are not equivalent*

(a)

$p$	$q$	$r$	$(p \Rightarrow q) \Rightarrow r$			$p \Rightarrow (q \Rightarrow r)$		
$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$
$T$	$T$	$F$	$T$	$F$	$F$	$T$	$F$	$F$
$T$	$F$	$T$	$F$	$T$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$F$	$T$	$F$	$T$	$T$	$T$
$F$	$T$	$T$	$T$	$T$	$T$	$F$	$T$	$T$
$F$	$T$	$F$	$T$	$F$	$F$	$F$	$T$	$F$
$F$	$F$	$T$	$T$	$T$	$T$	$F$	$T$	$T$
$F$	$F$	$F$	$T$	$F$	$F$	$F$	$T$	$T$

←

(b)  $p \Rightarrow (q \Rightarrow r)$  and  $(p \wedge q) \Rightarrow r$

*These formulas are equivalent; their truth tables are the same*

$p$	$q$	$r$	$p \Rightarrow (q \Rightarrow r)$			$(p \wedge q) \Rightarrow r$		
$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$
$T$	$T$	$F$	$T$	$F$	$F$	$T$	$F$	$F$
$T$	$F$	$T$	$T$	$T$	$T$	$F$	$T$	$T$
$T$	$F$	$F$	$T$	$T$	$T$	$F$	$T$	$F$
$F$	$T$	$T$	$F$	$T$	$T$	$F$	$T$	$T$
$F$	$T$	$F$	$F$	$F$	$F$	$F$	$T$	$F$
$F$	$F$	$T$	$F$	$T$	$T$	$F$	$T$	$T$
$F$	$F$	$F$	$F$	$F$	$F$	$F$	$T$	$F$

(c)  $(p \wedge q) \Rightarrow r$  and  $(p \Rightarrow r) \wedge (q \Rightarrow r)$

Since their truth tables differ, these formulas are not equivalent

$p$	$q$	$r$	$(p \wedge q) \Rightarrow r$			$(p \Rightarrow r) \wedge (q \Rightarrow r)$		
$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$	$T$
$T$	$T$	$F$	$T$	$F$	$F$	$F$	$F$	$F$
$T$	$F$	$T$	$F$	$T$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$F$	$T$	$F$	$F$	$F$	$T$
$F$	$T$	$T$	$F$	$T$	$T$	$T$	$T$	$T$
$F$	$T$	$F$	$F$	$T$	$F$	$T$	$F$	$F$
$F$	$F$	$T$	$F$	$T$	$T$	$T$	$T$	$T$
$F$	$F$	$F$	$F$	$T$	$F$	$T$	$T$	$T$



7. For each of the following propositions, give the DNF under the variable ordering  $\langle a, b, c \rangle$ .

(a)  $a \vee (\neg a \wedge \neg b)$

(b)  $a \Rightarrow (b \Leftrightarrow c)$

(c)  $(\neg b \wedge c) \wedge (\neg a \Rightarrow \neg c) \wedge (c \wedge (\neg b \vee \neg a))$

(d)  $(a \Rightarrow b) \Leftrightarrow (b \Rightarrow c)$

SOLUTION

(a) Variable  $c$  does not appear in this formula, so it may be excluded from the truth table.

$a$	$b$	$c$	$a$	$\vee$	$(\neg a \wedge \neg b)$	
$F$	$F$	$-$	$T$	$T$	$T$	$\leftarrow$
$F$	$T$	$-$	$F$	$F$	$F$	
$T$	$F$	$-$	$T$	$F$	$F$	$\leftarrow$
$T$	$T$	$-$	$T$	$F$	$F$	$\leftarrow$

However, a DNF over  $\langle a, b, c \rangle$  might be required to include all  $c$ -clauses. This “full” DNF is

$$(\neg a \wedge \neg b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c) \vee (\neg a \wedge b \wedge \neg c) \vee (\neg a \wedge b \wedge c) \vee (a \wedge \neg b \wedge \neg c) \vee (a \wedge \neg b \wedge c) \vee (a \wedge b \wedge \neg c) \vee (a \wedge b \wedge c)$$

If  $c$  is omitted, the resulting DNF is

$$(\neg a \wedge \neg b) \vee (\neg a \wedge b) \vee (a \wedge b)$$

(b)

$a$	$b$	$c$	$a \Rightarrow (b \Leftrightarrow c)$	
$F$	$F$	$F$	$T$	$\leftarrow$
$F$	$F$	$T$	$T$	$\leftarrow$
$F$	$T$	$F$	$T$	$\leftarrow$
$F$	$T$	$T$	$T$	$\leftarrow$
$T$	$F$	$F$	$T$	$\leftarrow$
$T$	$F$	$T$	$F$	
$T$	$T$	$F$	$F$	
$T$	$T$	$T$	$T$	$\leftarrow$

The DNF is

$$(\neg a \wedge \neg b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c) \vee (\neg a \wedge b \wedge \neg c) \vee (\neg a \wedge b \wedge c) \vee (a \wedge \neg b \wedge \neg c) \vee (a \wedge \neg b \wedge c)$$

(c) In the truth table below, working from left to right, I stopped once I reached a conjunct that evaluates to  $F$ . The **red** entries show where I stopped.

$a$	$b$	$c$	$(\neg b \wedge c)$	$\wedge$	$(\neg a \Rightarrow \neg c)$	$\wedge$	$(c \wedge (\neg b \vee \neg a))$
$F$	$F$	$F$	<b><math>F</math></b>	$F$	$T$	<b><math>F</math></b>	$F$
$F$	$F$	$T$	<b><math>T</math></b>	$F$	<b><math>F</math></b>	<b><math>F</math></b>	$T$
$F$	$T$	$F$	<b><math>F</math></b>	$F$	$T$	<b><math>F</math></b>	$F$
$F$	$T$	$T$	<b><math>F</math></b>	$F$	$F$	<b><math>F</math></b>	$T$
$T$	$F$	$F$	<b><math>F</math></b>	$F$	$T$	<b><math>F</math></b>	$F$
$T$	$F$	$T$	<b><math>T</math></b>	<b><math>T</math></b>	<b><math>T</math></b>	<b><math>T</math></b>	<b><math>T</math></b>
$T$	$T$	$F$	<b><math>F</math></b>	$F$	$T$	<b><math>F</math></b>	$F$
$T$	$T$	$T$	<b><math>F</math></b>	$F$	$T$	<b><math>F</math></b>	$F$

The DNF contains no clauses. The “or” of zero terms evaluates to  $F$ , so one could write  $DNF \equiv T$ .

(d)

$a$	$b$	$c$	$(a \Rightarrow b)$	$\Leftrightarrow$	$(b \Rightarrow c)$	
$F$	$F$	$F$	$T$	<b><math>T</math></b>	$T$	←
$F$	$F$	$T$	$T$	<b><math>T</math></b>	$T$	←
$F$	$T$	$F$	$T$	$F$	$F$	
$F$	$T$	$T$	$T$	<b><math>T</math></b>	$T$	←
$T$	$F$	$F$	$F$	$F$	$T$	
$T$	$F$	$T$	$F$	$F$	$T$	
$T$	$T$	$F$	$T$	$F$	$F$	
$T$	$T$	$T$	$T$	<b><math>T</math></b>	$T$	←

$$DNF \equiv (\neg a \wedge \neg b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c) \vee (\neg a \wedge b \wedge c) \vee (a \wedge b \wedge c)$$

8. Reduce the following boolean expressions to simpler terms

(a)  $xy + (x + y)\bar{z} + y$

(b)  $x + y + \overline{(x + y + z)}$

(c)  $yz + wx + z + [wz(xy + wz)]$

SOLUTION

(a)  $xy + (x + y)\bar{z} + y$   
 $= xy + y + (x + y)\bar{z}$  *commutativity*  
 $= xy + 1y + (x + y)\bar{z}$  *Identity*  
 $= yx + y1 + (x + y)\bar{z}$  *commutativity, twice*  
 $= y(x + 1) + y\bar{z} + x\bar{z}$  *distributivity*  
 $= y1 + y\bar{z} + x\bar{z}$  *dominance*  
 $= y(1 + \bar{z}) + x\bar{z}$  *distributivity*  
 $= y1 + x\bar{z}$  *dominance*  
 $= y + x\bar{z}$  *identity*

For the remaining derivations, rearrangement by commutativity and associativity is implicit.

(b)  $x + y + \overline{(x + y + z)}$   
 $= x + y + (\bar{x} \bar{y} \bar{z})$  *DeMorgan's Law*  
 $= x + y + (x \bar{y} \bar{z})$  *negation*  
 $= y + x + (x \bar{y} \bar{z})$  *commutativity*  
 $= y + x1 + x \bar{y} \bar{z}$  *identity*  
 $= y + x(1 + \bar{y} \bar{z})$  *distributivity*  
 $= y + x$  *dominance, identity*

(c)  $yz + wx + z + [wz(xy + wz)]$   
 $= yz + wx + z + wzxy + wz$  *distributivity, twice*  
 $= yz + wx + z + wzxy + wz$  *idempotence*  
 $= wx + z + yz + wzxy + wz$  *commutativity, associativity*  
 $= wx + z(1 + y + wxy + w)$  *distributivity*  
 $= wx + z$  *dominance, identity*

9. Write the truth tables for the following logical formulas and state whether each is a tautology, a contradiction, or neither (a contingency).

(a)  $P \wedge (Q \vee R)$

(b)  $(P \wedge \neg P) \Rightarrow Q$

(c)  $P \Rightarrow (Q \vee \neg Q)$

SOLUTION

$R$	$P$	$Q$	(a)			(b)			(c)		
			$P$	$\wedge$	$(Q \vee R)$	$(P \wedge \neg P)$	$\Rightarrow$	$Q$	$P$	$\Rightarrow$	$(Q \vee \neg Q)$
$T$	$T$	$T$	$T$	$T$	$T$	$F$	$T$	$T$	$T$	$T$	$T$
$T$	$T$	$F$	$T$	$T$	$T$	$F$	$T$	$F$	$T$	$T$	$T$
$T$	$F$	$T$	$T$	$T$	$T$	$F$	$T$	$T$	$F$	$T$	$T$
$T$	$F$	$F$	$T$	$F$	$F$	$F$	$T$	$F$	$F$	$T$	$T$
$F$	$T$	$T$	$F$	$F$	$T$	—	—		—	—	—
$F$	$T$	$F$	$F$	$F$	$F$	—	—		—	—	—
$F$	$F$	$T$	$F$	$F$	$T$	—	—		—	—	—
$F$	$F$	$F$	$F$	$F$	$F$	—	—		—	—	—
			<i>contingency</i>			<i>tautology</i>			<i>tautology</i>		

COMMENTS.

1. I chose to arrange the truth table variables in the order  $R, P, Q$  because the formula in Parts (b) and (c) contain only  $P$  and  $Q$ . This way, I could use just the upper half of the table as a truth table for (c).
2. In Part (b) once I recognized that the premise  $P \wedge \neg P$ , is a contradiction, I realized that evaluation of the implication  $(P \wedge \neg P) \Rightarrow Q$  “reduces” to  $F \Rightarrow Q$ , which is always  $T$ .
3. Similarly, in Part (c), once I recognized that the conclusion  $Q \vee \neg Q$  is always  $T$ , and  $P \Rightarrow T$  is always  $T$

No SUPPLEMENTAL PROBLEM this week.